## UNIVERSAL FAMILIES OF ARCS AND CURVES ON SURFACES

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## SURFACES

- A surface is a topological space that locally looks like a plane or a half plane.

- We deal with connected, compact, orientable surfaces determined by their genus and number of punctures.



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- A surface is a topological space that locally looks like a plane or a half plane.

- We deal with connected, compact, orientable surfaces determined by their genus and number of punctures.
- Closed curves are maps from $S^{1}$ to the surface. We deal with simple curves.



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- A surface $S_{g, b}$ is decomposed to $2 g-2+b$ pairs of pants.
$\rightarrow$ A pants decomposition of $S_{g, b}$ contains $3 g-3+b$ curves.


## REALIZING ALL PANTS DECOMPOSITIONS

- Two pants decompositions are of the same homeomorphism type if there exists a homeomorphism of the surface that maps one to the other.

- A dual trivalent graph that encodes the adjacencies of the pants determines the homeomorphism type of a pants decomposition.

- A pants universal family of curves is a set of curves that contains all types of pants decomposition of the surface


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- A pants universal family of curves is a set of curves that contains all types of pants decomposition of the surface
* Problem: What is the size of this family?


## Estimating The size of THE universal family

- A surface of genus $g$ has $g^{\Theta(g)}$ homeomorphism types of pants decompositions.
- There exists an upper bound of $g^{O(g)}$ for the size of such a family.


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- A surface of genus $g$ has $g^{\Theta(g)}$ homeomorphism types of pants decompositions.
- There exists an upper bound of $g^{O(g)}$ for the size of such a family.
- We can improve this upper bound by reusing the curves:


Theorem (F.,de Mesmay, Parlier)
Let $\Gamma$ be a minimal size universal family for pants decompositions of a surface of genus $g$ without punctures. Then

$$
|\Gamma| \leq 3^{2 g-1} \quad \text { and } \quad|\Gamma|=\Omega\left(g^{\frac{4}{3}-\epsilon}\right) \quad \text { for any } \quad \epsilon>0 .
$$

Question: Can we close the gap between the upper and lower bound?

## REALIZING PANTS DECOMPOSITIONS

Given a sphere with $n$ punctures, we can consider two kinds of homeomorphisms:

- homeomorphisms that fix punctures pointwise (labelled punctures)
- homeomorphisms that can permute the punctures (unlabelled punctures)


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## A MOTIVATION: UNIVERSAL SHORTEST PATH METRIC

## Theorem (FÀRY-WAGNER)

All planar graphs can be drawn in the plane with straight lines.

## Conjecture (Hubard, Kaluža, Tancer) '17

For any surface $S_{g}$, there exists a universal Riemannian metric (universal shortest path metric such that all graphs embeddable on $S_{g}$ can be embedded with shortest paths.

- It is true for sphere, the projective plane, torus and the Klein bottle.
- For higher $g$, most metrics are not a universal shortest path metric.


## Theorem (Hubard, Kaluža, de Mesmay, Tancer '17)

For every $g>1$, there exists a universal metric for which all graphs embeddable on the surface can be embedded such that edges are concatenations of $O(g)$ shortest paths.

## A motivation: Universal shortest path metric

## Observation

A super-polynomial lower bound on the size of a pants universal family of curves leads to a negation of the universal shortest metric conjecture.

- For any pants decomposition of the surface create a graph embedding that contains the decomposition
- Assuming the conjecture, all these graphs have a shortest path embedding
$\rightarrow$ each curve in the pants decomposition embeds with constant number of shortest paths

$\rightarrow$ Each pair of curves in the family cross constantly


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$\rightarrow$ Each pair of curves in the family cross constantly
- Maximal number of non-homotopic simple closed curves on a surface that cross constantly is polynomial in $g$. [Greene '18] $\%$


## The case of surfaces Without punctures

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$\rightarrow$ Upper bound:

- Any pants decomposition contains $g$ simple closed curves that cut the surface into a sphere with $2 g$ punctures.


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- The size of this family is less than $3^{n-1}$.


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$\rightarrow$ Lower bound:

- There exists up to $g^{g}$ types of pants decompositions.
- Each pants decomposition consists of $3 g-3$ curves.

$$
\binom{|\Gamma|}{3 g-3} \gtrsim g^{g} \rightarrow|\Gamma|=\Omega\left(g^{\frac{4}{3}-\epsilon}\right)
$$

Question: Can we do any better?

## THE CASE OF SPHERE WITH UNLABELLED PUNCTURES

## Theorem

Let $\Gamma$ be a pants universal family of curves on with minimal size on the sphere with $n$ (unlabelled) punctures. Then

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|\Gamma|=O\left(n^{2}\right) \quad \text { and } \quad|\Gamma|=\Omega(n \log (n))
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$\rightarrow$ Lower bound:


■ For $i \leq \frac{n}{2}$ at least $\left\lfloor\frac{n}{i}\right\rfloor$ curves that separate $i$ boundaries must exist in the pants universal family.
$\rightarrow \Sigma\left\lfloor\frac{n}{i}\right\rfloor=\Omega(n \log n)$.

## Families of edges realizing all Triangulations

## REALIZING ALL TRIANGULATIONS

- Two triangulations of an n-gon are of the same type, if one can be obtained from a rotation and/or reflection of the other.

- This is similar to the case of spheres with (unlabelled) boundaries.



## REALIZING ALL TRIANGULATIONS

## Theorem (F., de Mesmay, Parlier)

Let $E$ be a family of edges that realize all the traingulations of an n-gon.
We have $|E|=O\left(n^{2}\right)$ and $|E|=\Omega\left(n^{2-\epsilon}\right) \quad$ for any $\epsilon>0$.
$\rightarrow$ Upper bound follows by computing all the diagonals in the $n$-gon.
$\rightarrow$ Lower bound:
■ Any such family also needs to realize all types of m-cycles.

- There exist $\Theta\left(n^{m-1}\right)$ types of $m$-cycles.
- Any such family defines a graph with $|E|$ edges.
$(2,1,2,1)$
- any graph with $|E|$ edges has $O\left(|E|^{\frac{m}{2}}\right)$ $m$-cycles. (Rivin)



## REALIZING ALL TYPES OF PANTS/ TRIANGLES

- A simpler question is to ask for the size of a family of curves that realize all types of triangles.

- or on a sphere with $n$ (unlabelled) punctures, a family that realizes all types of pairs of pants:



## Summary and conclusion

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## Thank You!

